# On vertex covers, matchings and random trees

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#### Abstract

We pay special attention to the corresponding backbones i.e. these vertices that are occupied and those that are empty in every minimal vertex cover (resp. these egdes that are occupied and those that are empty in every maximal matching). The key result in our approach is that for trees, the backbones can be recovered from a particular tri-coloring which has a simple characterization. We give applications to the computation of some averages related to the enumeration of minimal vertex covers and maximal matchings in the random labeled tree ensemble, both for finite size and in the asymptotic regime.

## 1 Motivations

For a given (simple: no loops, no multiple edges) graph, finding the size of a minimal vertex cover or of a maximal matching (see the beginning of sec.2 for a reminder of definitions), and counting the number of solutions all fall

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under the generic name of combinatorial optimization problems, a field with a long history.

To analyse the average behavior for these questions, the simplest model is the Erdös-Renyi model of random graphs. In this context, the evaluation of the average size of a maximal matching has been solved by Karp and Sipser [1] (see [2] for refinements and [3] for a physicist approach) in the thermodynamic limit, i.e. when both V, the vertex set, and  $\mathcal{E}$ , the edge set become large, but the ratio  $\alpha = 2|\mathcal{E}|/|V|$  has a finite limit. For the average size of a minimal vertex cover, the answer is known only when  $\alpha \leq e$  and asymptotically for large  $\alpha$  [4].

To get more detailed informations on these problems, one can investigate several combinatorial patterns. In this paper, we concentrate on backbones, i.e. these vertices that are occupied and those that are empty in every minimal vertex cover (resp. these egdes that are occupied and those that are empty in every maximal matching). While for general graphs the relationships between the backbones are complicated, this is not true for trees: in that case the backbone geometry can be recovered from a special tricoloring, unique for each tree and which is easily characterized. This is the content of Theorem 1, the crucial ingredient for our subsequent analysis. We use this theorem to compute the average size of the backbones and the average number of minimal vertex covers and maximal matchings for random labeled trees of size n, where each of the  $n^{n-2}$  labeled trees has the same probability. We also analyse the asymptotic behavior for large n, see Theorems 2,4 and 5

Due to the simple, locally treelike, structure of Erdös-Renyi random graphs<sup>1</sup>, insight can often be obtained from an analysis of trees, even if in the case of this paper the extension is nontrivial. The present study of backbones and the corresponding applications to random labeled trees can thus be seen as a preliminary step towards the analysis of their random graph generalizations.

Our motivation to study backbones comes from physics. The adjacency matrix of a random graph (which is symmetric) can be seen as an example of a random Hamiltonian, whose average spectrum one would like to compute.

<sup>&</sup>lt;sup>1</sup>For  $\alpha \leq 1$ , an Erdös-Renyi random graph is a forest for most thermodynamical purposes, but the local treelike structure remains even after the birth of the giant component. Moreover, the finite components of size n are distributed with the uniform measure on labeled trees of size n in the thermodynamic limit.

In the infinite  $\alpha$  limit, one recovers a semi circle, but for finite  $\alpha$  matters are much more complicated. The spectrum contains a dense family of delta peaks plus presumably a continuous component when  $\alpha$  exceeds a threshold. In the case of the zero eigenvalue, the structure of the eigenvectors can be studied in detail<sup>2</sup>. It exhibits interesting phenomena of localization and delocalization when  $\alpha$  varies [6]. We shall see below that these phase transitions are closely related to the structure of the backbones. However, their combinatorial interpretation is still unclear to us.

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## 2 Main results

A vertex cover of the graph  $A = (V, \mathcal{E})$  is a subset of V containing at least one end of each edge in  $\mathcal{E}$ . We are interested in minimal vertex covers, i.e. those whose cardinality is the smallest. The positive (resp. negative) vertex-backbone of A is the set of vertices which belong to every (resp. no) minimal vertex cover. The other vertices are called degenerate vertices. An edge between degenerate vertices is called exclusive if no minimal vertex cover contains its two extremities.

A matching of A is a set of non-adjacent edges of A. Maximal matchings are those whose cardinality is the largest. The positive (resp. negative) edge-backbone is the set of edges which belong to every (resp. no) maximal matching. The other edges are called degenerate edges. A vertex for which there is a maximal matching none of whose edges contains it as an extremity is called optional. The vertices that are neither optional, nor an extremity of an edge in the positive backbone are called unavoidable.

If A is a tree or a forest, one can characterize these objects by simple properties and compute them recursively. If the  $n^{n-2}$  labeled trees on n vertices are chosen at random with the counting measure, we shall use this to adress questions of the type "What is the average size of the edge-backbones?" or "What is the average number of maximal vertex covers" in the random labeled tree ensemble.

Note that we are interested in global extrema. A local version for, say,

<sup>&</sup>lt;sup>2</sup>This is indeed an example of a situation where the analysis of random trees proved crucial to understand the case of the Erdös-Renyi model.

minimal vertex covers would be vertex covers such that changing the state of any occupied vertex to the empty state destroys the vertex cover property<sup>3</sup>. Some problems analogous to the ones we deal with but for local problems can be found for instance in the work of Meir and Moon, see e.g. [7] and references therein. For local extrema problems, the notion of backbones seems to be less relevant.

A tricoloring of the graph  $A = (V, \mathcal{E})$  is a triple  $(B, \mathcal{R}, G) \subset V \times \mathcal{E} \times V$ , such that B,G and the set of end-vertices of  $\mathcal{R}$  form a partition of V. As a starting point,

**Theorem 1** Suppose A is a tree. Each of the three properties (i), (ii) and (iii) characterizes one and the same tricoloring  $(B, \mathcal{R}, G)$  of A.

- (i) Minimal vertex-covers: B is the positive backbone;  $\mathcal{R}$  is the set of exclusive edges; G is the negative backbone.
- (ii) Maximal matchings: B is the set of unavoidable vertices;  $\mathcal{R}$  is the positive backbone; G is the set of optional vertices.
- (iii) The edges in  $\mathcal{R}$  are non-adjacent; the edges with one end-vertex in G have the other end-vertex in B; each vertex in B is connected to G by at least two edges.

This unique tricoloring  $(B, \mathcal{R}, G)$  is called the *b-coloring* of A. An edge is said red if it is in  $\mathcal{R}$ , and a vertex is said brown if it lies in B, green if it lies in G and red if it is an end-vertex of a red edge.

In the sequel,  $N_c(A)$  denotes the number of vertices with color c (where c is either brown, red or green) in the tree A and  $N_c(n)$  is the total number of vertices with this same color among the  $n^{n-2}$  labeled trees on n vertices. We shall work with the generating functions  $F_c(x) \equiv \sum_{n\geq 1} \frac{N_c(n)}{n!} x^n$ . They all involve the tree generating function  $T(x) \equiv \sum_{n\geq 1} \frac{n^{n-1}}{n!} x^n$ . Our main combinatorial and probabilistic results are contained in the following theorems.

**Theorem 2** The generating functions for the total number of brown, red and green vertices are

$$F_B = T(x) + T(-T(x)) - T(-T(x))^2$$
;  $F_R = T(-T(x))^2$ ;  $F_G = -T(-T(x))$ 

and the corresponding explicit first terms, closed formulæ, and asymptotics

$$N_B = (0, 0, 3, 4, 185, 1026, 30457, 362664, 10245825, 195060070, \cdots)$$

<sup>&</sup>lt;sup>3</sup>The example of a starlike tree shows the difference between the local and global versions.

$$\frac{N_B(n)}{n^{n-1}} = 1 + \sum_{l=1}^n \left(\frac{-l}{n}\right)^l \left(\frac{2}{l} - 1\right) \binom{n}{l} \sim 0.2276096757 \cdots 
N_R = (0, 2, 0, 48, 120, 4560, 35700, 1048992, 15514128, 456726240, \cdots) 
\frac{N_R(n)}{n^{n-1}} = -2 \sum_{l=1}^n \left(\frac{-l}{n}\right)^l \left(\frac{1}{l} - 1\right) \binom{n}{l} \sim 0.4104940676 \cdots 
N_G = (1, 0, 6, 12, 320, 2190, 51492, 685496, 17286768, 348213690, \cdots) 
\frac{N_G(n)}{n^{n-1}} = -\sum_{l=1}^n \left(\frac{-l}{n}\right)^l \binom{n}{l} \sim 0.3618962567 \cdots$$

Corollary 3 The size (that is, the cardinality) of the minimal vertex covers and maximal matchings of a tree A is  $N_B(A) + N_R(A)/2$ , hence the average fraction of vertices in a vertex cover of a tree on n vertices is  $n^{1-n}(N_B(n) + N_R(n)/2) \sim 0.4328567095$  for large n.

Let  $N_{vc}(n)$  and  $N_m(n)$  denote the total numbers of minimal vertex covers and of maximal matchings among labeled trees on n vertices. The corresponding generating functions,  $F_{vc}(x) \equiv \sum_{n\geq 1} \frac{N_{vc}(n)}{n!} x^n$  and  $F_m(x) \equiv \sum_{n\geq 1} \frac{N_m(n)}{n!} x^n$ , verify

**Theorem 4** The generating function for the total number of minimal vertex covers is

$$F_{vc}(x) = (1 - U)xe^{U} - UT(x^{2}e^{2U}) + U - \frac{1}{2}U^{2},$$

where  $xUe^U = T(x^2e^{2U})(e^{xe^U} - 1)$ .  $N_{vc} = (1, 2, 3, 40, 185, 3936, 35917, 978160, 14301513, 464105440, \cdots)$ 

**Theorem 5** The generating function for the number of maximal matchings is

$$F_m(x) = -\frac{1}{2}(xe^U + U)^2 + (1 + Uxe^U)xe^U + U - U^2,$$

where  $U = x^2 e^{-x^2 e^{2U} + xe^U + 3U}$ .

 $N_m = (1, 1, 6, 24, 320, 3270, 55482, 999656, 21718440, 544829130, \cdots)$ 

**Remark 6**. Sketch of the relation with the kernel of the adjacency matrix (see [6] for details).

The kernel of the adjacency matrix of a tree is directly related to the b-coloring. First one shows, for instance by induction on the size of the tree, that the kernel of A has dimension  $N_G(A) - N_B(A)$ . Second, one shows that the support <sup>4</sup> of the kernel consists of the green vertices.

Moreover, the maximal subsets B', G' of V such that

(iii)' the edges with one end-vertex in G' have the other end-vertex in B' and each vertex in B' is connected to G' by at least two edges,

coincide with B and G of the b-coloring of A. Thus, in the case of trees, maximality allows to define the sets B and G without mentionning  $\mathcal{R}$ .

Drawing the edges between B' and G' defines a bicolored subforest of A. But there is a partial converse to these constructions: one can show that, for a general graph, a bicolored subforest on B', G' satisfying (iii)' allows to define a |G'| - |B'| dimensional subspace of the kernel with support G'. For the Erdös-Renyi model, the enumeration of the finite maximal bicolored subtrees satisfying (iii)' accounts for the full dimension of the kernel up to an o(|V|) correction for small or large  $\alpha$ , but there is a window of  $\alpha$ 's for which infinite patterns contribute O(|V|) to the dimension of the kernel. These are the localization-delocalization transitions alluded to before.

These are the results that motivated us to have a closer look at the backbones.

## 3 Proof of theorem 1

There are many ways to build a proof, depending on personal tastes, and the choice of the authors has been subject to many fluctuations. So it is not unlikely that the reader will spare time finding his own argument instead of understanding the proof that we propose.

Let  $A = (V, \mathcal{E})$  be a tree. Property (ii) of theorem 1 obviously characterizes a unique tricoloring of A. Hence, it suffices to establish that

- Step 1 :  $(B, \mathcal{R}, G)$  defined by (i) is a tricoloring, and it satisfies (iii);
- Step 2:  $(iii) \Rightarrow (ii)$ .

 $<sup>^4</sup>$ The vectors on which the adjacency matrix acts can be interpreted as maps from V to the reals and it makes sense to say that a vector vanishes on a given vertex. The support of a vector is the set of vertices on which it does not vanish. The support of a family of vectors is the union of the elementary supports.

If  $A = (\{v\}, \emptyset)$  is the isolated vertex, any of the three assertions (i, ii, iii) defines a unique tricoloring  $(\emptyset, \emptyset, \{v\}) - v$  is green – and we suppose from now on that A has at least two vertices.

## 3.1 Step 1

Let  $(B, \mathcal{R}, G)$  be the triple defined by (i) in theorem 1. Because B, G and the set of end-vertices of  $\mathcal{R}$  are mutually disjoint, proving that a degenerate vertex is the end of an exclusive edge should ensure that  $(B, \mathcal{R}, G)$  is a tricoloring of A. Then we check that it satisfies (iii).

Deletion of  $v \in V$  and its incident edges leaves  $p \geq 1$  trees  $A_1, \dots, A_p$ , with  $A_i = (V_i, \mathcal{E}_i)$ . Denote by  $v_i$  the unique vertex of  $A_i$  which is adjacent to v in A.

A vertex cover of A obviously induces a vertex cover on each  $A_i$ . Conversely, suppose we are given a vertex cover  $C_i$  on each  $A_i$ , and denote by C the union of the  $C_i$ 's. An edge of A is either in some  $\mathcal{E}_i$ , in which case it has one end in  $C_i$  hence in C, or one of the p edges between v and some  $v_i$ . Hence  $C \cup \{v\}$  is a vertex cover of A, but C is not unless it contains each of the  $v_i$ 's.

Now, let us write  $n_i$  (resp.  $\bar{n}_i$ ) for the minimal cardinality of vertex covers of  $A_i$  containing (resp. not containing)  $v_i$ . As a consequence of the previous remarks, a subset C of V is a minimal vertex cover of A if and only if one of two exclusive assertions holds:

- Assertion  $1: 1 + \sum_{i} \min(n_i, \bar{n}_i) \leq \sum_{i} n_i$ , C contains v and C induces a minimal cover on each  $A_i$ .
- Assertion  $2: 1 + \sum_{i} \min(n_i, \bar{n}_i) \geq \sum_{i} n_i$ , C does not contain v and C induces on each  $A_i$  a vertex cover of cardinality  $n_i$  containing  $v_i$ .

This gives us constraints for v being or not in some backbone, which are very informative if we note that  $n_i \leq \bar{n}_i + 1$  for all i.

Suppose first that v is degenerate. Then  $1 + \sum_{i} \min(n_i, \bar{n}_i) = \sum_{i} n_i$ , and this implies that  $n_{i_0} = \bar{n}_{i_0} + 1$  for a unique  $i_0$ . There exists a minimal vertex cover of A containing v, which induces a minimal vertex cover on  $A_{i_0}$ , hence does not contain  $v_{i_0}$ . There exists also a minimal vertex cover not containing v, which obviously contains  $v_{i_0}$ : as was to be proved,  $v_{i_0}$  is degenerate, and v is the end of an exclusive edge  $\{v, v_{i_0}\}$ .  $(B, \mathcal{R}, G)$  is thus a tricoloring. Now, given  $i \neq i_0$ , we can find a minimal vertex cover of  $A_i$  containing  $v_i$  because

 $n_i \leq \bar{n}_i$  and then extend it into a minimal vertex cover of A containing both v and  $v_i$ . Hence v is actually the end of a unique exclusive edge, proving that the edges of  $\mathcal{R}$  are not adjacent.

If v is in the negative backbone, a minimal vertex cover of A does not contain v, hence it contains each  $v_i$ . So the neighbors of vertices in G are in B.

If v is in the positive backbone,  $1 + \sum_i \min(n_i, \bar{n}_i) < \sum_i n_i$ , which proves the existence of at least two distinct i's such that  $n_i = \bar{n}_i + 1$ . The corresponding  $A_i$ 's do not admit minimal covers containing  $v_i$ . Since a minimal cover of A contains v, it induces a minimal cover on these  $A_i$ 's: it does not contain the (at least two) corresponding  $v_i$ 's. Hence every vertex in B has at least two neighbors in G.

This proves that  $(B, \mathcal{R}, G)$  in (i) is a tricoloring which satisfies (iii) and we now come to the proof of  $(iii) \Rightarrow (ii)$ .

### 3.2 Step 2

Let  $(B, \mathcal{R}, G)$  be a tricoloring of A satisfying (iii). Let R denote the set of end-vertices of edges in  $\mathcal{R}$ .

If  $B = G = \emptyset$ , then  $\mathcal{R}$  is a perfect matching of the tree A. Because a tree admits at most one perfect matching  $^5$ ,  $\mathcal{R}$  is the unique maximal matching of A.

Relax this assumption, suppose  $e_0 = \{g_0, b_0\}$  is an edge of A ( $g_0 \in G, b_0 \in B$ ) and let  $\mathcal{M}$  be a matching of A containing  $e_0$ . Obviously, there exist paths of the form  $g_0, b_0, \dots, g_k, b_k$  ( $k \geq 0$ ) such that  $\{g_i, b_i\} \in \mathcal{M}, g_i \in G$  and  $b_i \in B$  for all i. Take one with maximal length. Then  $b_k$  has at least one neighbor  $g_{k+1} \in G \setminus \{g_k\}$ , which is not the end of an edge in  $\mathcal{M}$  (because our path is maximal, and all edges ending at  $g_{k+1}$  have the other end in  $g_0$ ). Now, we replace in  $g_0$ 0 the  $g_0$ 1 to  $g_0$ 2. This leads to a matching with same cardinality as  $g_0$ 3, but not containing the edge  $g_0$ 3.

As a first consequence, there exist maximal maximal matchings not containing  $g \in G$  as an end-vertex (vertices in G are optional). Now, let  $b \in B$  and suppose  $\mathcal{M}$  is a matching not containing b as an end-vertex. We show that  $\mathcal{M}$  is not maximal. If some neighbor v of b is not the end-vertex of any

<sup>&</sup>lt;sup>5</sup>This is clear by induction on the size of the tree, if we note that an edge ending at a leaf is contained in any perfect matching.

edge in  $\mathcal{M}$  we can append the edge  $\{b, v\}$  to  $\mathcal{M}$ . On the other hand, suppose that some neighbor of  $b, g_0 \in G$ , is the end of an edge in  $\mathcal{M}$ . Then, we apply the procedure above to build a matching  $\mathcal{M}'$  of A with same cardinality as that of  $\mathcal{M}$ , not containing  $g_0$  as an end-vertex. But these two matchings coincide except on some path  $g_0, b_0, \dots, g_k, b_k, g_{k+1}$ , which does not contain  $b \neq b_0$ . Hence,  $\mathcal{M}'$  does not contain any edge ending at b or  $g_0$ , and we can append the edge  $\{b, g_0\}$ . Hence vertices in B are unavoidable.

Anticipating the conclusion of this pragraph, let us call forbidden edges the edges between two vertices in B, between a vertex in B and one in R, and those edges not in R with both ends in R. Deletion of the forbidden edges leaves some trees, and  $(B, \mathcal{R}, G)$  induces on each of these trees a tricoloring satisfying (iii), of the form  $(\emptyset, \mathcal{R} \cap \mathcal{E}_i, \emptyset)$  or  $(B \cap V_i, \emptyset, G \cap V_i)$ . Moreover, a matching  $\mathcal{M}$  induces a matching on each of these trees and, if  $\mathcal{M}$  contains  $p \geq 1$  forbidden edges, at least p+1 of these induced matchings do not contain some vertex in B as an end-vertex or some edge in  $\mathcal{R}$ . By the preceding remarks, they are not maximal. Hence, by deleting from  $\mathcal{M}$  these p edges, and by replacing the edges of these p+1 matchings by those of maximal matchings, we obtain a matching of A containing at least one more edge than  $\mathcal{M}$ . So a maximal matching does not contain any forbidden edge and, as an easy corollary, deletion of these edges leaves maximal matchings of the resulting trees.

Thus, edges in  $\mathcal{R}$  are in the positive backbone of A. Denote by  $(\hat{B}, \hat{\mathcal{R}}, \hat{G})$  the tricoloring of A defined by (ii): we have proved that  $B \subset \hat{B}, \mathcal{R} \subset \hat{\mathcal{R}}, G \subset \hat{G}$ . By general properties of tricolorings, this implies that  $(B, \mathcal{R}, G) = (\hat{B}, \hat{\mathcal{R}}, \hat{G})$  and concludes the proof of theorem 1.

**Remark** Let us consider a minimal vertex cover C of A. It contains all the  $N_B(A)$  brown vertices of A and none of the green vertices. The other  $N_R(A)$  vertices are ends of non-adjacent red edges, and we have seen that C contains exactly one end of each such edge: hence C contains exactly  $N_B(A) + N_R(A)/2$  vertices.

A maximal matching of A contains all the  $N_R(A)/2$  red edges and exactly one edge ending at each brown vertex, the other end being green (hence not brown). Moreover it does not contain any other edge: hence a maximal matching of A contains exactly  $N_B(A) + N_R(A)/2$  edges. This proves corollary 3.

#### 4 Generating functions

#### Generating function for b-colorings 4.1

Our purpose in this section is to give an exponential generating function for the number of labeled trees with given color distribution:

$$F(g, b, r) \equiv \sum_{n>1} \sum_{A \in \mathcal{A}_n} \frac{1}{n!} g^{N_G(A)} b^{N_B(A)} r^{N_R(A)}$$

where  $A_n$  is the set of labeled trees on n vertices.

Recalling that a tree has a unique b-coloring, we say that a rooted tree has color c if its root has color c. Let G, B, R be the (exponential) generating functions for respectively green, brown, red rooted trees.

Let A be a rooted tree. Then

• A is green if, and only if, its root is connected to the root of arbitrarily many trees defined as follows: root adjacent to arbitrarily many rooted colored trees, with the condition that at least one root be green. Let us call quasi-brown these trees and denote by U their generating function. Then

$$G = ge^{U}$$

$$U = be^{B+R}(e^{G} - 1)$$

$$(1)$$

$$(2)$$

$$U = be^{B+R}(e^G - 1) (2)$$

• A is brown if, and only if, its root is connected to the root of arbitrarily many brown or red rooted trees and to at least two green rooted trees, SO

$$B = be^{B+R}(e^G - 1 - G) (3)$$

• Finally, A is red if, and only if, its root is connected to arbitrarily many brown or red rooted trees and to exactly one tree defined as follows: root adjacent to arbitrarily many red or brown rooted trees. Let us call quasi-red these trees and denote by Q their generating function. Then

$$R = rQe^{B+R}$$

$$Q = re^{B+R}$$

$$(5)$$

$$Q = re^{B+R} (5)$$

Now, the generating function F for colored trees is the only function of g, b, r such that  $g\frac{\partial F}{\partial g} = G$ ,  $b\frac{\partial F}{\partial b} = B$ ,  $r\frac{\partial F}{\partial r} = R$  and F(0,0,0) = 0. The following F indeed satisfies these conditions, thus the generating function for colored trees is

$$F(g,b,r) = -\frac{1}{2}((B+R)^2 + Q^2) - GU + be^{B+R}(e^G - 1 - G) + ge^U + rQe^{B+R}$$
 (6)

We check that F(x, x, x) gives back the usual generating function for labeled trees:  $F_0(x) = \sum_{n \geq 1} \frac{n^{n-2}}{n!} x^n$ . Recall that the generating function  $T = xF_0'(x)$  for rooted trees verifies  $T(x) = xe^{T(x)}$  for |x| < 1/e. Putting S = B + R and taking g, b, r = x in the equations (3)+(4)-(2) and (5)-(1) yields

$$S - U = (Q - G)xe^{S}$$

$$Q - G = xe^{U}(1 - e^{S-U})$$

Taking  $x \to 0$ , this implies S = U and G = Q, so (2)+(5) yields  $S + G = xe^{S+G}$ . Hence S + G = S + Q = T(x), and it follows from (5) that  $Qe^Q = xe^{S+Q} = xe^{T(x)} = T(x)$ . Finally:

$$S = U = T(x) + T(-T(x))$$
  
$$G = Q = -T(-T(x))$$

Inject these into F(x, x, x) to get  $F = T(x) - \frac{1}{2}T(x)^2$ , which is indeed equal to  $F_0(x)$ .

Remark 7. Sketch of the relation with Feynman graph enumeration.

Define S(S = B + R, U, G, Q, g, b, r) by the right hand-side of eq.(6) but seen as a function of seven independent variables. Then the vanishing of the partial derivative of S with respect to S leads to the combination eq.(3)+eq.(4), whereas the vanishing of the partial derivatives with respect to U, G and Q leads to eq.(1), eq.(2) and eq.(5) respectively. Thus, F is the value of S at the (unique in the small g, b, r expansion) extremum in the capital variables. The same kind of considerations would apply to all the generating functions in this paper

We do not know if there is a simple combinatorial explanation for this extremal property, but there is a simple physical interpretation that we give

in appendix A to illustrate how two scientific communities deal with the same problem. The reader interested in a more thourough study of the combinatorics of Feynman graphs can consult e.g. [10].

The generating function for the total number of vertices with a given color comes from differentiation with respect to the corresponding variable, followed by the identification b=g=r=x:

$$\sum_{n} \frac{N_{B}(n)}{n!} x^{n} = T(x) + T(-T(x)) - T(-T(x))^{2}$$

$$\sum_{n} \frac{N_{G}(n)}{n!} x^{n} = -T(-T(x))$$

$$\sum_{n} \frac{N_{R}(n)}{n!} x^{n} = T(-T(x))^{2}$$

In order to give explicit formulæ for the average numbers of vertices of each color, we need to know the term of given degree in T(-T(x)) and  $T(-T(x))^2$ . Writing these as contours integrals along a small contour surrounding 0 and changing the integration variable x into  $-te^t$  yields

$$\oint \frac{dx}{x^{n+1}} T(-T(x)) = \frac{(-1)^n}{n} \oint \frac{dt}{t^n} e^{-nt} T'(t) 
\oint \frac{dx}{x^{n+1}} T(-T(x))^2 = 2 \frac{(-1)^n}{n} \left( \oint \frac{dt}{t^n} e^{-nt} T'(t) - \oint \frac{dt}{t^{n+1}} e^{-nt} T(t) \right),$$

from which follow both the closed forms and, by the steepest descent method, their large-size asymptotics<sup>6</sup>

$$\begin{split} \frac{N_B(n)}{n^{n-1}} &= 1 + \sum_{l=1}^n \left(\frac{-l}{n}\right)^l \left(\frac{2}{l} - 1\right) \binom{n}{l} \sim 1 + T'(-1) + 2T(-1)) \\ \frac{N_G(n)}{n^{n-1}} &= -\sum_{l=1}^n \left(\frac{-l}{n}\right)^l \binom{n}{l} \sim T'(-1) \\ \frac{N_R(n)}{n^{n-1}} &= -2\sum_{l=1}^n \left(\frac{-l}{n}\right)^l \left(\frac{1}{l} - 1\right) \binom{n}{l} \sim -2(T'(-1) + T(-1)) \end{split}$$

<sup>&</sup>lt;sup>6</sup>In this simple situation, we can proceed naively to get the asymptotics. For a more rigorous treatment in a similar but slightly more involded context, see e.g. [8].

## 4.2 Generating function for minimal vertex covers

Let us define a covered tree to be a pair (A, C), where A is a rooted tree and C is a minimal vertex cover of A. The generating functions for brown and green covered trees are denoted respectively by B and G in this section. For red covered trees, it is useful to make the distinction between the minimal covers which contain the root and those which do not: let us denote by  $R_+, R_-$  the corresponding generating functions.

Consider a pair (A, C), where A is a rooted tree with root v and C a subset of the set of vertices of A. Then

- (A, C) is a green covered tree if, and only if,  $v \notin C$ , v is attached to arbitrarily many quasi-brown trees, and C induces on each of these trees a vertex cover with minimal cardinality among those containing the root.
- (A, C) is a brown covered tree if, and only if,  $v \in C$ , v is attached to at least 2 green trees and to arbitrarily many brown or red rooted trees, and C induces on each of these trees a minimal vertex cover.
- (A, C) is a red covered tree if, and only if, v is attached to exactly one quasi-red tree  $A_{i_0}$  and to arbitrarily many brown or red trees, and one of two exclusive assertions holds: (1)  $v \in C$ , v induces a minimal cover on each of the attached tree; (2)  $v \notin C$ , v induces on each of the attached trees a cover with minimal cardinality among those containing the root.

This leads to

$$B = b(e^{G} - 1 - G)e^{B+R_{+}+R_{-}} \qquad G = ge^{U}$$

$$R_{+} = rQ_{-}e^{B+R_{+}+R_{-}} \qquad R_{-} = rQ_{+}e^{B+R_{+}}$$

where the auxiliary function  $U, Q_+, Q_-$  are defined as

$$U \equiv b(e^G - 1)e^{B + R_+ + R_-}, \quad Q_+ \equiv re^{B + R_- + R_+}, \quad Q_- \equiv re^{B + R_+}.$$

The generating function for covered rooted trees is  $G + B + R_+ + R_-$  and

$$F_{vc} = ge^{U} + b(e^{G} - 1 - G)e^{B+R_{+}+R_{-}} + rQ_{-}e^{B+R_{+}+R_{-}} + rQ_{+}e^{B+R_{+}} - GU - \frac{1}{2}(B^{2} + R_{+}^{2}) - R_{+}R_{-} - B(R_{+} + R_{-}) - Q_{+}Q_{-}$$

turns out to be the only function with correct b, r, g partial derivatives satisfying  $F_{vc}(0, 0, 0) = 0$ .

Let us identify b, r, g = x. Then  $B + R_+ = U$ ,  $G = Q_-$  and  $R_+ = R_- \equiv R = x^2 e^{2U+R} = T(x^2 e^{2U})$ . Hence the closed formula for U is  $xUe^U = (e^{xe^U} - 1)T(x^2 e^{2U})$ , and the expression for  $F_{vc}$  follows immediately.

## 4.3 Generating function for maximal matchings

We shall skip the details, the crucial points being that

- A maximal matching of a tree A contains all the red edges and exactly one edge ending at each brown vertex, the other end being green. It does not contain any other edge.
- Given an edge B-G, there exist maximal matchings which do not contain it (because g is optional). There also exist some which do contain it. Indeed, let  $e=\{b,g\}$  be an edge of A with  $b\in B, g\in G$ . There exists a maximal matching not containing g as an end vertex and, as a maximal one, this matching contains an edge e' ending at b. Just replace e' by e.

Then the generating functions for matched trees read

$$G_{+} = gU_{-}e^{U_{+}} \qquad G_{-} = ge^{U_{+}}$$
 
$$B = bG_{-}(e^{G_{+}+G_{-}} - 1)e^{B+R} \qquad R = rQe^{B+R}$$

where

$$U_{+} \equiv bG_{-}e^{G_{+}+G_{-}+B+R}, \quad U_{-} \equiv b(e^{G_{+}+G_{-}}-1)e^{B+R}, \quad Q \equiv re^{B+R}$$

The generating function writes

$$F_m = gU_-e^{U_+} + ge^{U_+} + bG_-(e^{G_++G_-} - 1)e^{B+R} + rQe^{B+R} -G_+U_+ - G_-(U_+ + U_-) - \frac{1}{2}(B+R)^2 - \frac{1}{2}Q^2.$$

For b=r=g=x, we find  $B+R=U_+\equiv U, Q=G_-$ . Hence,  $G_+=U-x^2e^{2U}$  and  $U=x^2e^{3U+xe^U-x^2e^{2U}}$ , as was to be proved.

**Remark 8.** The quantities  $\frac{1}{n} \log N_{vc}(n)$ ,  $\frac{1}{n} \log N_m(n)$  given analytically by theorems 4,5 are difficult to confront to numerical simulations sampling the

 $n^{n-2}$  trees uniformly. They are typical examples of non self-averaging quantities. This means basically that a small fraction of trees contributes significantly to the average although it is unlikely to be "visited" in reasonable time by a Monte-Carlo algorithm sampling trees uniformly. A simpler task is the numerical estimation of  $\frac{1}{n} < \log N_{vc}(A) >$  or  $\frac{1}{n} < \log N_m(A) >$ , a self averaging quantity which answers the question: how many minimal vertex covers or maximal matchings does a typical tree have. This question will be adressed analytically and numerically in a work to come, again based on the use of b-colorings [9].

## A A note on Feynman graphs

In quantum field theory, graph counting occurs for the following reasons. A physical system is characterized by an action  $\mathcal{S}(T_a, g_i)$  where the  $T_a$ 's denote dynamical variables and  $g_i$  coupling constants. The index a often runs through a continuum, but for the present dicussion, we assume it to take a finite number of values. This is the usual case that there is only a finite number of coupling constants. The general structure of  $\mathcal{S}$  is  $\mathcal{S} = \mathcal{S}_0 + \sum_i g_i P_i(T_a)$  where  $\mathcal{S}_0 = -\frac{1}{2} \sum_{a,b} C_{ab} T_a T_b$  is a quadratic form which we assume here to be nondegenerate and the  $P_i(T_a)$  are analytic functions. The quantity to be computed is the free energy

$$\hbar \log \int \prod_{a} \frac{dT_a}{\sqrt{2\pi\hbar}} \sqrt{\det C} \exp \frac{\mathcal{S}(T_a, g_i)}{\hbar},$$
(7)

where the integration contours and values of the  $g_i$ 's are choosen to ensure convergence of the integral, and  $\hbar$  is Plank's constant. Note that when the coupling constants  $g_i$  all vanish, this expression vanishes too. The so-called semi-classical expansion expresses the free energy of the system as an asymptotic expansion in powers of  $\hbar$ , the  $\hbar^n$  term being computable by definite rules from certain non simple connected graphs, the so-called Feynman graphs, with n independent cycles. To understand the appearance of graphs, the easy way is to first expand formally the integrand in (7) in powers of the coupling constants and then the  $P_i$ 's in powers of the fields  $T_a$ . This reduces the integral to integration of a monomial against a gaussian weight, and the combinatorics of the result is obtained by repeated integration by parts, which amounts to pair successively and in all possible ways all pairs

of variables  $T_aT_b$  in the monomial and replace them by  $\hbar(C^{-1})_{ab}$ . In that way, one can interpret the monomial in the  $P_i$ 's as vertices marked with the fields they involve, and  $\hbar(C^{-1})_{ab}$  is the weight for an edge of type ab between two vertices: all possible graphs that can be built in that way appear in the formal power series expansion of the integral. Taking the logarithm to compute the free energy amounts to keep only connected graphs as usual in combinatorics.

To make this general idea concrete take  $S = -T^2/2 + ge^T$ . In that case the quadratic form is associated to the  $1 \times 1$  identity matrix and  $e^T$  describes vertices of arbitrary degree. The integral along the real axis with a purely imaginary g makes sense and on can compute the asymptotic expansion at small g. The result is that

$$\int \frac{dT}{\sqrt{2\pi\hbar}} \exp \frac{\mathcal{S}(T,g)}{\hbar} \sim \sum_{n>0} \frac{(g\hbar^{-1})^n}{n!} e^{n^2\hbar/2}$$

On the other hand, if to an arbitrary graph (not necessarily simple) on n vertices, described by a symmetric matrix  $M = (m_{pq})$ , one gives a weight

$$\prod_{p < q} \frac{\hbar^{m_{pq}}}{m_{pq}!} \prod_{p} \frac{1}{2^{m_{pp}}}$$

which is essentially its symmetry factor, the sum over all graphs reconstructs the factor  $e^{n^2\hbar/2}$ .

The classical limit corresponds to keeping only the  $\hbar^0$  contribution, i.e. trees. On the other hand, in the classical limit the system is decribed by the classical equations of motion, which say that the action  $\mathcal{S}$  is extremal with respect to all field variations: in the present context, this boils down to the stationary phase method. And indeed, for the concrete example above, this extremum condition leads to  $T = ge^T$ , so that T is the rooted labeled tree generating function.

We leave it to the reader to compute the inverse of C, to extract the kinds of vertices and edges that Feynman graphs produce when

$$S(T, U, G, Q, g, b, r) = -\frac{1}{2}(T^2 + Q^2) - GU + be^T(e^G - 1 - G) + ge^U + rQe^T$$

and retrieve in that way the conditions (iii).

The main message is that, while for a combinatorist eqs.(1,2,3,4,5) for rooted trees follow from routine arguments, for a quantum field theorist it is

S(T, U, G, Q, g, b, r) which comes immediately to mind to count the desired unrooted graphs and, at the extremum in (T, U, G, Q), trees.

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